

# A criterion for internality of some differential equations

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# Parametrizing solutions of differential equations

Consider two complex numbers  $a$  and  $b$ , and the very basic differential equations:

- ▶  $y' = a$

- ▶  $y' = by$

the solution sets can be parametrized as:

- ▶  $\{y_0 + c, c \in \mathbb{C}\}$

- ▶  $\{cy_1, c \in \mathbb{C}\}$

where  $y_0, y_1$  are particular solutions of the equations. Or even:

- ▶  $\{at + c, c \in \mathbb{C}\}$

- ▶  $\{ce^{bt}, c \in \mathbb{C}\}$

## Systems of equations

$$\begin{cases} y_1' = y_1 \\ y_2' = iy_2 \\ z' = 6 \end{cases} \Rightarrow \begin{cases} y_1 \in \{c_1 e^t, c_1 \in \mathbb{C}\} \\ y_2 \in \{c_2 e^{it}, c_2 \in \mathbb{C}\} \\ z \in \{6t + d, d \in \mathbb{C}\} \end{cases}$$

Any three solutions are algebraically independent, so we get three independent parametrizations. Three particular solutions,  $e^t$ ,  $e^{it}$  and  $t$ , are needed.

$$\begin{cases} y_1' = 2y_1 \\ y_2' = 4y_2 \\ z_1' = 3 \\ z_2' = 6 \end{cases} \Rightarrow \begin{cases} y_1 \in \{c_1 e^{2t}, c_1 \in \mathbb{C}\} \\ y_2 \in \{c_2 e^{4t}, c_2 \in \mathbb{C}\} \Rightarrow y_2 = \frac{c_2}{c_1^2} y_1^2 \\ z_1 \in \{3t + d_1, d_1 \in \mathbb{C}\} \\ z_2 \in \{6t + d_2, d_2 \in \mathbb{C}\} \Rightarrow z_2 = 2z_1 + d_2 - 2d_1 \end{cases}$$

Only two functions,  $e^t$  and  $t$ , are needed to parametrize the set of solutions.

## More complicated example

$$\begin{cases} y' = \frac{yz}{y+z} \\ z' = -\frac{yz}{y+z} \end{cases} \Rightarrow \begin{cases} y \in \left\{ \frac{ce^x}{e^x - d}, c \in \mathbb{C}, d \in \mathbb{C}^* \right\} \\ z = c - y \end{cases}$$

So we can still parametrize the solutions as rational functions of  $e^x$ , but not in a linear way. This is because this system is in (non-linear) bijection with a linear system:

$$\begin{cases} y' = \frac{yz}{y+z} \\ z' = -\frac{yz}{y+z} \end{cases} \Rightarrow \begin{cases} u = \frac{y}{z} \\ v = y + z \end{cases} \Rightarrow \begin{cases} u' = u \\ v' = 0 \end{cases}$$

## Our goal

Consider some system of differential equations of the general form:

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_n(y_1, \dots, y_n) \end{cases}$$

where  $f_i \in \mathbb{C}(x_1, \dots, x_n)$ .

*What we saw:* it is possible for solutions of such a system to be in rational bijection with solutions of a linear system. In that case, we obtain a rational parametrization by "transferring" the linear one.

*What we'll do:* the converse is true! If such a system has a parametrization using rational functions, then it must be in rational bijection with a linear system.

## How a model theorist thinks about this

We want a convenient structure to work with differential equations.

Bare minimum: differential fields of characteristic zero, i.e. fields equipped with a differential  $\delta$  that is additive and satisfies Leibniz's rule  $\delta(ab) = \delta(a)b + a\delta(b)$ . We will often denote  $\delta(a) = a'$ .

The theory of differential fields of characteristic zero has a model companion, which is the theory  $\text{DCF}_0$  of differentially closed fields.

Concretely, this means that if  $K \models \text{DCF}_0$  and some finite system of differential (in)equations, defined over some parameters  $A \subset K$ , has a solution in some differential field extension  $K < L$ , then it has a solution in  $K$ .

$\text{DCF}_0$  has **quantifier elimination**: any formula is equivalent to a boolean combination of differential equation and inequations.

## Types in $\text{DCF}_0$

It will be more convenient to work with **types**, instead of definable sets. Fix some  $M \models \text{DCF}_0$ .

### Definition

Given  $b \in M$  and  $A \subset M$ , the type of  $b$  over  $A$  is the set of all formulas, with parameters in  $A$ , that are satisfied by  $b$ . We denote it  $\text{tp}(b/A)$ .

In general, a type  $p$  over  $A$  is a maximal consistent set of formulas with parameters in  $A$ . We let  $S(A)$  be the set of types over  $A$ .

By quantifier elimination,  $\text{tp}(b/A)$  is just the set of differential equations and inequations, with parameters in  $A$ , satisfied by  $b$ .

## An important example

We will care about systems of the form:

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_n(y_1, \dots, y_n) \end{cases}$$

where the  $f_i$  are rational functions.

Such a system has a **generic type**  $p$ : this is the type of some  $a_1, \dots, a_n$  satisfying these equations, but no other non-trivial differential equation.

Because the differential equations have order one, this means simply no polynomial equation.



## Fixing some model

For the rest of the talk, I will fix some monster model  $\mathcal{U} \models \text{DCF}_0$ , or equivalently, a model that is homogeneous and saturated in its own cardinality.

Concretely, if  $A \subset \mathcal{U}$  and  $|A| < |\mathcal{U}|$ , then for any  $p \in S(A)$ :

- ▶ the set of **realizations** of  $p$  in  $\mathcal{U}$ , i.e. elements of  $\mathcal{U}$  satisfying all formulas in  $p$ , is non-empty. We denote it  $p(\mathcal{U})$ , and write  $a \models p$  for  $a \in p(\mathcal{U})$ .
- ▶ if  $b \in p(\mathcal{U})$ , then  $p(\mathcal{U})$  is the orbit of  $b$  under the action of  $\text{Aut}(\mathcal{U}/A)$ .

To define what I mean by "parametrizing", I need the  $\emptyset$ -definable set of **constants**:

$$\mathcal{C} = \{x \in \mathcal{U}, \delta(x) = 0\}$$

Thinking  $\mathcal{C} = \mathbb{C}$  is fine.

# Rational parametrization = Internality

$F$  will always be an algebraically closed differential field.

## Definition

A type  $p \in S(F)$  is  $\mathcal{C}$ -internal if there are

- ▶  $a_1, \dots, a_n$  realizations of  $p$
- ▶ an  $F$ -definable function  $f(x_1, \dots, x_n, y_1, \dots, y_m)$

such that for all  $a \models p$ , there are  $c_1, \dots, c_m \in \mathcal{C}$  with:

$$a = f(a_1, \dots, a_n, c_1, \dots, c_m)$$

If we replace  $f$  with a one-to-finite correspondence, then we say  $p$  is almost  $\mathcal{C}$ -internal.

# What we want

## Question

Is there a criterion to determine whether a type  $p \in S(F)$  is almost  $\mathcal{C}$ -internal?

We think of types as representing generic solutions of a differential equations. We will examine systems of the form:

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_n(y_1, \dots, y_n) \end{cases}$$

where the  $f_i \in F(x_1, \dots, x_n)$  and  $F$  is a field of *constants*.

A general system may have some polynomial equations between the  $y_i$ . We will not deal with this more general case.

# Binding groups

Internal types are structured by the following theorem:

## Theorem

*If  $p \in S(F)$  is  $\mathcal{C}$ -internal, then the group action of restrictions to  $p(\mathcal{U})$  of automorphisms of  $\mathcal{U}$  is isomorphic to an  $F$ -definable group action.*

*It is called the **binding group of  $p$** , and denoted  $\text{Aut}_F(p/\mathcal{C})$ .*

Moreover, the group  $\text{Aut}_F(p/\mathcal{C})$  is definably isomorphic to  $G(\mathcal{C})$ , for some algebraic group  $G$ .

Key properties:

- ▶ if  $\text{Aut}_F(p/\mathcal{C})$  acts transitively, we say  $p$  is **weakly  $\mathcal{C}$ -orthogonal**.
- ▶ if  $\text{Aut}_F(p/\mathcal{C})$  acts freely (i.e. without fixed point), we say  $p$  is **fundamental**.

## Weakly orthogonal and fundamental

Fact (Kolchin, model-theoretic translation by Jaoui-Moosa)

*Let  $p \in S(F)$  be a  $\mathcal{C}$ -internal, weakly  $\mathcal{C}$ -orthogonal and fundamental type. Then there is an algebraic group  $G$  defined over  $F \cap \mathcal{C}$  such that  $p$  is interdefinable (i.e. in  $F$ -definable bijection) with the generic type  $q$  of the solution to a full logarithmic differential equation on  $G$  over  $F$ .*

$\text{Aut}_F(p/\mathcal{C})$  must be definably isomorphic to  $G(\mathcal{C})$

What we can do:

- (A) reduce to weakly  $\mathcal{C}$ -orthogonal and fundamental types
- (B) control what  $G$  can appear as a binding group
- (C) write concrete equations for the solution to a full logarithmic differential equation
- (D) use interdefinability to obtain an explicit condition for internality

## (B) Linear binding groups are abelian

We will only need the two most basic algebraic groups:

- ▶  $G_a(\mathcal{C}) = (\mathcal{C}, +)$ ,
- ▶  $G_m(\mathcal{C}) = (\mathcal{C} \setminus \{0\}, \cdot)$ .

### Fact

*Let  $F$  be a field of constants and  $p \in S(F)$  be an internal, weakly  $\mathcal{C}$ -orthogonal type. If  $\text{Aut}_F(p/\mathcal{C})$  is linear, then it is isomorphic to  $G_m(\mathcal{C})^k \times G_a(\mathcal{C})^l$ , where  $k \in \mathbb{N}$  and  $l \in \{0, 1\}$ .*

The action of the binding group is always faithful, and a faithful transitive action of an abelian group is always free, i.e.  $p$  must be fundamental!

## (B) The binding group is linear

Consider  $p$  the generic type of some system:

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_n(y_1, \dots, y_n) \end{cases}$$

where the  $f_i \in F(x_1, \dots, x_n)$  and  $F$  is a field of constants.

We see that:

- ▶ the action of  $\text{Aut}_F(p/\mathcal{C})$  is definably isomorphic to some birational action of an algebraic group  $G(\mathcal{C})$  on the affine space  $\mathbb{A}^n(\mathcal{C})$
- ▶ algebraic geometry  $\Rightarrow$  the binding group is linear

By the previous slide, if  $p$  is weakly  $\mathcal{C}$ -orthogonal, then  $\text{Aut}_F(p/\mathcal{C})$  is definably isomorphic to  $G_m(\mathcal{C})^k \times G_a(\mathcal{C})^l$  for some  $k \in \mathbb{N}$  and  $l \in \{0, 1\}$ .

## (C) Logarithmic differential equations

*To summarize:* let  $F$  be a field of constants, and  $p \in S(F)$  be a  $\mathcal{C}$ -internal, weakly  $\mathcal{C}$ -orthogonal type. Then  $p$  is interdefinable with the generic type of a full logarithmic differential equation on  $(G_m)^k \times (G_a)^l$ , with  $k \in \mathbb{N}$  and  $l \in \{0, 1\}$ .

Such an equation can be expressed by (if  $l = 1$ ):

$$\begin{cases} z'_1 = \lambda_1 z_1 \\ \vdots \\ z'_k = \lambda_k z_k \\ z'_{k+1} = 1 \end{cases}$$

and fullness is equivalent to the  $\lambda_i$  being  $\mathbb{Q}$ -linearly independent.

A dimension argument shows that it's either  $(G_m)^{n-1} \times G_a$  or  $(G_m)^n$ , i.e.  $k + 1 = n$



## (D) What interdefinability gives

Assume we are in the  $(G_m)^{n-1} \times G_a$  case.

$$\begin{cases} y'_1 = f_1(y_1, \dots, y_n) \\ \vdots \\ y'_n = f_n(y_1, \dots, y_n) \end{cases} \xrightarrow{F\text{-definable bijection}} \begin{cases} z'_1 = \lambda_1 z_1 \\ \vdots \\ z'_{n-1} = \lambda_{n-1} z_{n-1} \\ z'_n = 1 \end{cases}$$

By quantifier elimination:

the definable bijection is given by rational maps

$g_1, \dots, g_{n-1}, h \in F(x_1, \dots, x_n)$  such that:

$$\begin{cases} g_1(y_1, \dots, y_n)' = \lambda_1 g_1(y_1, \dots, y_n) \\ \vdots \\ g_{n-1}(y_1, \dots, y_n)' = \lambda_{n-1} g_{n-1}(y_1, \dots, y_n) \\ h(y_1, \dots, y_n)' = 1 \end{cases}$$

## Main theorem in the weakly orthogonal case

### Theorem (Eagles-J.)

Let  $F$  be an algebraically closed field of constants, some  $f_1, \dots, f_n \in F(x_1, \dots, x_n)$  and  $p$  the generic type of the system:

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_n(y_1, \dots, y_n) \end{cases}$$

Then  $p$  is almost  $\mathcal{C}$ -internal and weakly  $\mathcal{C}$ -orthogonal if and only if there are rational functions  $g_1, \dots, g_{n-1}, h \in F(x_1, \dots, x_n)$ ,  $\mathbb{Q}$ -linearly independent  $\lambda_1, \dots, \lambda_{n-1} \in F$  with:

- ▶  $\sum_{i=1}^n \frac{\partial g_j}{\partial x_i} f_i = \lambda_j g_j$  for all  $1 \leq j \leq n-1$ ,
- ▶  $\sum_{i=1}^n \frac{\partial h}{\partial x_i} f_i = 1$ .

or some similar equations for the  $(G_m)^n$  case.

## What about the non-weakly $\mathcal{C}$ -orthogonal case?

### Example

The generic type of:

$$\begin{cases} y' = \frac{yz}{y+z} \\ z' = -\frac{yz}{y+z} \end{cases}$$

is internal, and not weakly  $\mathcal{C}$ -orthogonal:

$(y+z)' = 0$ , so  $y+z \in \mathcal{C}$ , which must be fixed by the binding group  $\Rightarrow$  the binding group does not act transitively.

Non-weak  $\mathcal{C}$ -orthogonality was witnessed by a definable function  $(y, z) \rightarrow y+z$  to  $\mathcal{C}$ . This is true in general:

### Lemma

*Let  $F$  be an algebraically closed field of constants and  $p \in S(F)$ . Then there is an  $F$ -definable map  $\pi : p \rightarrow \pi(p)$  such that  $\pi(p)(\mathcal{U}) \subset \mathcal{C}$  and for any  $a \models p$ , the type  $\text{tp}(a/\pi(a)F)$  is stationary and weakly  $\mathcal{C}$ -orthogonal.*

## (A) Reducing to weakly $\mathcal{C}$ -orthogonal: proof idea

Let  $F$  be an algebraically closed field of constants and some internal type  $p \in S(F)$ .

Consider the map  $\pi$  from the previous slide, so for any  $a \models p$ :

- ▶  $\pi(a) \in \mathcal{C}$
- ▶  $\text{tp}(a/\pi(a)F)$  is stationary and weakly  $\mathcal{C}$ -orthogonal

Then:

- ▶  $\text{Aut}_F(p/\mathcal{C})$  is linear  $\Rightarrow \text{Aut}_{\pi(a)F}(\text{tp}(a/\pi(a)F)/\mathcal{C})$  is also linear
- ▶  $\text{tp}(a/\pi(a)F)$  is a type over constant parameters. We can (modulo technicalities) apply our previous theorem to get rational maps to logarithmic-differential equations on  $G_m$  or  $G_a$
- ▶  $\pi$  is also given by rational maps, we can pick them to be algebraically independent elements of  $F(x_1, \dots, x_n)$

## Theorem (Eagles-J.)

The generic type  $p$  of the system:

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_k(y_1, \dots, y_n) \end{cases}$$

is almost  $\mathcal{C}$ -internal if and only if there are rational functions  $g_1, \dots, g_{k-1}, h, h_1, \dots, h_{n-k} \in F(x_1, \dots, x_n)$ ,  $\mathbb{Q}$ -linearly independent  $\lambda_1, \dots, \lambda_{k-1} \in F$  with:

- ▶  $\sum_{i=1}^n \frac{\partial g_j}{\partial x_i} f_i = d_j g_j$  for all  $1 \leq j \leq k-1$  and  $\sum_{i=1}^n \frac{\partial h}{\partial x_i} f_i = 1$ ,
- ▶ the  $h_j$  are algebraically independent over  $F$  and  $\sum_{i=1}^n \frac{\partial h_j}{\partial x_i} f_i = 0$  for all  $1 \leq j \leq n-k$ .

or some similar equations if the binding group is  $(G_m)^k$ .

## Orthogonality to the constants

What if they are no rational functions  $g_j$  or  $h_j$ ? This corresponds to the model-theoretic notion of **orthogonality to the constants**:

### Theorem (Eagles-J.)

*The generic type  $p$  of the system:*

$$\begin{cases} y_1' = f_1(y_1, \dots, y_n) \\ \vdots \\ y_n' = f_n(y_1, \dots, y_n) \end{cases}$$

*is orthogonal to the constants if and only if there are no*

*$g \in F(x_1, \dots, x_n)$  and  $\gamma \in F$  such that  $\sum_{i=1}^n \frac{\partial g}{\partial x_i} f_i = \gamma g$  or  $= \gamma$ .*

Orthogonality to the constants implies that the generic solutions  $f_i$  are not **Liouvillian**: they cannot be constructed using elementary functions, composition, and integration.

## An application: the classic Lotka-Volterra system

The Lotka-Volterra system models predator-prey populations:

- ▶  $x$  represents the prey population,
- ▶  $y$  represents the predator population,

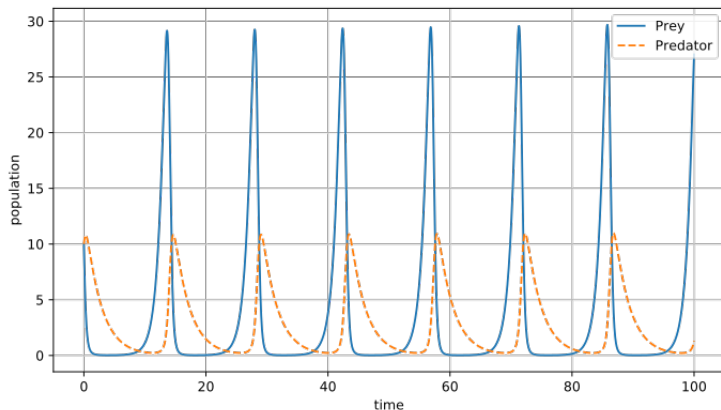
and is given by:

$$\begin{cases} x' = ax - bxy \\ y' = -cy + dxy \end{cases}$$

For realism, one may ask that  $a, b, c, d$  are strictly positive real numbers. We instead pick  $a, b, c, d \in \mathcal{C} \setminus \{0\}$ .

## Graphs of (real) solutions

$$\begin{cases} x' = ax - bxy & \text{prey} \\ y' = -cy + dxy & \text{predator} \end{cases}$$



Credit: Ian Alexander (parameters, PNG version) Krishnavedala (vectorisation), from wikipedia.



## Mostly not Liouvillian

### Theorem (Eagles-J.)

*Unless  $a = c$ , the generic solution of the Lotka-Volterra system:*

$$\begin{cases} x' = ax + bxy \\ y' = cy + dxy \end{cases}$$

*is not Liouvillian. If  $a = c$  it is elementary (proved by Varma [3]).*

Using our theorem, it is enough to show that the partial differential equations (where  $\mu = \frac{a}{c}$ ):

$$c \frac{\partial g}{\partial x_0} (\mu x_0 + x_0 x_1) + c \frac{\partial g}{\partial x_1} (x_1 + x_0 x_1) = \begin{cases} 0 \\ 1 \\ \lambda g \quad (\lambda \in \mathbb{Q}(a, b, c, d)^{\text{alg}}) \end{cases}$$

have no rational solutions. We use Laurent series.

## Further work

- ▶ add polynomial equations. Issue: The binding group need not be linear anymore. But the Chevalley decomposition should help.
- ▶ work over non constant parameters.
  - ▶ probably not as nice of a result: any algebraic group can appear as a binding group.
  - ▶ hope in low dimension. The case  $n = 1$  has essentially been solved by Jaoui-Moosa [2]. If  $n = 2$ , we are essentially interested in connected algebraic groups acting on  $\mathbb{P}^2$ , which were classified by Enriques [1].
- ▶ can model theory say anything about parametrizations by non-rational functions? For example solutions of  $y''y - (y')^2 = 0$  are  $\{ce^{dx} : c, d \in \mathbb{C}\}$ . The generic type is not almost  $\mathcal{C}$ -internal, essentially because  $x \rightarrow e^x$  is not definable in  $\text{DCF}_0$ .

Thank you!

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